# A polynomial interpolation algorithm for estimating a numerical function 

## Romulus Militaru


#### Abstract

We consider $(n+1)$ distinct numbers $x_{0}, x_{1}, \ldots, x_{n}$ and the values of the function $f$ at these $(n+1)$ given points. We study the problem of estimating $f(x)$ at a point $x=x^{*}$, within a given accuracy $\varepsilon$, using polynomial interpolation. Based on an analysis of the linear interpolation and the error of interpolation, we present an algorithm which better exploits the results of the calculus, contributing to a decrease in the amount of work involved, respectively the computational cost necessary to approximate $f\left(x^{*}\right)$ with a given error. The material also presents numerical examples solved using this algorithm.


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## 1. Introduction

Regardless of wheather we are talking about experimental results or about tables, most of the functions are given only for a discrete set of points which make up a sample. Obviously the cases in which the value of the function is present in the table for the point that we need are seldom, and hence the necessity to call for the interpolation, which is a basic tool for the approximation of a given function.

Consider that we are given the values of the function $f$ at $n+1$ distinct points $x_{0}, x_{1}, \ldots, x_{n}$. These are not necessarily equally spaced or even arranged in increasing order, the only main restriction is that they must be distinct. Occasionally the values of derivatives of $f$ are also prescribed.

The polynomial interpolation, which is included in the principal class of the linear interpolation, is the basis of several types of numerical integration formulas and is used in the construction of extrapolation methods for integration and differential equations, (see [2], [3], [8]).

The polynomial interpolation problem consists of determining a polynomial $P_{n}$ which takes the same values as $f$ at the $n+1$ points $x_{i}, i=0,1, \ldots, n$, (see [5], [2]). We could write that for the $n+1$ given pairs of number $\left(x_{k}, f_{k}\right), k=0,1, \ldots, n$ with $x_{i} \neq x_{j}$ for $i \neq j$

$$
P_{n}\left(x_{k}\right)=f_{k}, \quad k=0,1, \ldots, n
$$

## 2. Problem statement

We consider the problem of estimating $f(x)$ at a point $x=x^{*}$, within a given accuracy, using polynomial interpolation at distinct points $x_{i}, i=\overline{0, n}$, with a minimal computational cost.

The main problem is that estimating the size of the error $f-P_{n}$ from a knowledge of the values of $f$ at $x_{0}, x_{1}, \ldots, x_{n}$ alone, would be highly unlikely. Further information about $f$, such as bounds of his derivatives, are necessary, (see [7], [3]).
In consequence the degree of the polynomial needed for estimating $f\left(x^{*}\right)$ within a given accuracy $\varepsilon$ is generally not known.
We will consider in the sequel the cases of the Lagrange, Newton and Hermite interpolating polynomial. Concerning the accuracy of polynomial interpolation, there is the following result, (see [2], [5]):
Proposition 2.1. Let $[a, b]$ be any interval which contains all $n+1$ points $x_{0}, x_{1}, \ldots, x_{n}$. Let $f, f^{\prime}, \ldots, f^{(n)}$ exist and be continuous on $[a, b]$ and let $f^{(n+1)}$ exists for $x \in(a, b)$. Then, given any $x \in[a, b]$ there exist a number $\xi=\xi(x)$ in $(a, b)$ so that

$$
\begin{equation*}
f(x)-P(x)=\frac{f^{(n+1)}(\xi)}{(n+1)!} \omega(x) \tag{1}
\end{equation*}
$$

where $\omega(x)=\prod_{i=0}^{n}\left(x-x_{i}\right), P_{n}$ denoting the interpolating polynomial at the points $x_{i}, i=\overline{0, n}$.

Corollary 2.1. Under the hypotheses of Proposition 2.1, if we denote by $E(x)$ the error of interpolation, then

$$
\begin{equation*}
\|E\|_{\infty} \leq \frac{1}{(n+1)!}\left\|f^{(n+1)}\right\|_{\infty}\|\omega\|_{\infty} \tag{2}
\end{equation*}
$$

where $\|\phi\|_{\infty}=\sup \{|\phi(t)|, t \in[a, b]\}, \quad(\forall) \phi:[a, b] \rightarrow \mathbf{R}$ continue.

One can minimize the right side of (2) by choosing the interpolating points $x_{i}, i=\overline{0, n}$ as the zeros of the Cebyshev polynomials, (see [2],[5]),

$$
T_{n+1}:(-1,1) \rightarrow \mathbf{R}, \quad T_{n+1}(z)=\frac{1}{2^{n}} \cos [(n+1) \arccos (z)]
$$

Of course, this can be done if we have access to the values of $f$ at these zeros.
Remark 2.1. If the interpolating polynomial $P_{n}(x)$ is used to approximate $f(x)$ on the interval $[-1,1]$ then one can chooses $x_{i}=\cos \frac{\pi(2 i+1)}{2(n+1)}, i=0,1, \ldots, n$.
If $P_{n}(x)$ is used to approximate $f(x)$ on the interval $[a, b]$ then one can chooses the interpolating points given by $x_{i}=\frac{a+b}{2}+\frac{b-a}{2} \cos \frac{\pi(2 i+1)}{2(n+1)}, i=0,1, \ldots, n$.

Remark 2.2. In practice the estimating problem of a function $f(x)$ at a point $x=x^{*}$ presents itself mostly in the form: what is desired is the value $f\left(x^{*}\right)$ approximated with a given accuracy $\varepsilon$. Based on the above results, this tolerance $\varepsilon$ should not be selected smaller than $\frac{\left\|f^{(n+1)}\right\|_{\infty}}{(n+1)!}\|\omega\|_{\infty}$ with respect to the zeros of the Cebyshev polynomials

Let $H_{n}$ be the Hermite interpolating polynomial which satisfies

$$
\left\{\begin{array}{l}
d^{0}\left(H_{n}\right) \leq m-1 \\
H_{n}^{(k)}=f^{(k)}\left(x_{i}\right), 0 \leq i \leq n, 0 \leq k \leq \alpha_{i}
\end{array}\right.
$$

where

$$
\left\{\begin{array}{l}
x_{i}, i=\overline{0, n} \text { are the interpolating points, } \\
\alpha_{i}, i=\overline{0, n} \text { are positive integers, } \\
m=\sum_{i=0}^{n}\left(\alpha_{i}+1\right)
\end{array}\right.
$$

$H_{n}$ is the polynomial of lowest degree with the property that it agrees with the function $f$ and with all its derivatives of order less than or equal to $\alpha_{i}$ at $x_{i}, i=\overline{0, n}$.

Remark 2.3. If $\alpha_{i}=0, i=\overline{0, n}, H_{n}(x)$ becomes the Lagrange polynomial which interpolates the function $f$ at the points $x_{i}, i=\overline{0, n}$.

Denoting by $E(x)=f(x)-H_{n}(x), x \in[a, b]-$ an intervall which contains all points $x_{i}, i=\overline{0, n}$. Similar to Proposition 2.1, we have:

$$
\begin{equation*}
E(x)=\frac{f^{(m)}(\xi)}{m!} \omega_{m}(x), \quad \xi=\xi(x),(\forall) x \in[a, b] \tag{3}
\end{equation*}
$$

where $\omega_{m}(x)=\prod_{i=0}^{n}\left(x-x_{i}\right)^{\alpha_{i}+1}$ and $f \in C^{m}[a, b]$.
Under some restrictive conditions one states the following result:
Proposition 2.2. Consider that $f \in C^{m}[a, b]$ is an analytical one in the neighbourhood of $h=\frac{a+b}{2}$ with a convergence radius given by $r_{c}>\frac{3}{2}(b-a)$. Then the following statement holds:

$$
\|E\|_{\infty} \leq A \cdot\left(\frac{b-a}{p}\right)^{m} \cdot\left(1+\left(2^{m+1}-1\right) \frac{b-a}{2 p-b+a}\right) \quad(\forall) p \in\left(\frac{b-a}{2}, r_{c}\right)
$$

where $A=A(p)$ is a constant, and the sequence $\left(H_{n}\right)_{n \geq 1}$ of Hermite interpolating polynomials converges uniformly to the function $f$.
Proof. $f(x)=\sum_{i \geq 0} a_{i}(x-h)^{i}, \quad(\forall) x \in\left(h-r_{c}, h+r_{c}\right)$
Choosing $p \in\left(\frac{b-a}{2}, r_{c}\right), x=h+p$ then $\sum_{i \geq 0} a_{i} p^{i}$ converges.
In particular $\lim _{i \rightarrow \infty} a_{i} p^{i}=0$. There is a constant $A=A(p)$ so that $\left|a_{i}\right| \leq \frac{A}{p^{i}}, i \geq 0$.
Thus it follows that

$$
\left|f^{(m)}(x)\right| \leq A \cdot \sum_{i \geq m} \frac{1}{p^{i}} \frac{i!}{(i-m)!}|x-h|^{i-m}
$$

$(\forall) x \in[a, b] \subset\left(h-r_{c}, h+r_{c}\right)$ so that $|x-h| \in\left[0, \frac{b-a}{2}\right]$
Based on (3) one obtains

$$
\begin{equation*}
\|E\|_{\infty} \leq \frac{(b-a)^{m}}{m!}\left\|f^{(m)}\right\|_{\infty} \tag{4}
\end{equation*}
$$

Taking into consideration that $(\forall) z \in\left[0, \frac{b-a}{2}\right] \subset[0, p)$
$\sum_{i \geq m} \frac{1}{p^{i}} \frac{i!}{(i-m)!} z^{i-m}=\left(\frac{1}{1-\frac{z}{p}}\right)^{(m)} \leq \frac{m!p}{\left(p-\frac{b-a}{2}\right)^{m+1}}=\frac{m!}{p^{m}}\left(1+\frac{b-a}{2 p-b+a}\right)^{m+1}$ and using (4) one gets

$$
\|E\|_{\infty} \leq A \cdot \frac{(b-a)^{m}}{p^{m}} \cdot\left(\frac{2 p}{2 p-b+a}\right)^{m+1}=A \cdot \frac{p}{p-\frac{b-a}{2}} \cdot\left(\frac{b-a}{p-\frac{b-a}{2}}\right)^{m}
$$

If $p \geq \frac{3}{2}(b-a) \Leftrightarrow r_{c} \geq \frac{3}{2}(b-a)$ then $\left\|f-H_{n}\right\|_{\infty} \rightarrow 0$ uniformly when $m \rightarrow \infty$.
Using the inequality $(1+\beta)^{n} \leq 1+\left(2^{n}-1\right) \beta,(\forall) n \geq 1, \beta \in[0,1]$ it follows that

$$
\|E(x)\|_{\infty} \leq A \cdot\left(\frac{b-a}{p}\right)^{m} \cdot\left(1+\left(2^{m+1}-1\right) \frac{b-a}{2 p-b+a}\right)
$$

## 3. Choice of interpolating points

Taking into consideration that the amount of work necessary to approximate $f(x)$ at a point $x=x^{*}$ within a given accuracy $\varepsilon$, by means of the polynomial interpolation is proportional to the degree of interpolating polynomial used, in order to have a minimal computational cost one proposes the following technique:

Denote by $P_{k}(x)$ the interpolating polynomial of degree $\leq k$ which interpolates the function $f$ at distinct points $x_{i}, i=\overline{0, k}$. We calculate successively $P_{i}\left(x^{*}\right), i=0,1, \ldots$, increasing the number of interpolation points and hence the degree of the interpolating polynomial, until the difference between two consecutive values becomes smaller than $\varepsilon$, and thus a satisfactory approximation $P_{q}\left(x^{*}\right)$ to $f\left(x^{*}\right)$ has been found.
The main problem is how to choose the interpolating points for the successive construction of the interpolating polynomials $P_{i}(x), i=0,1, \ldots$.
In order to minimize the computational cost we start with the points $x_{i}, x_{i+1}$ so that $x^{*} \in\left(x_{i}, x_{i+1}\right)$, considering that $z_{0}=x_{i}, z_{1}=x_{i+1}$ if $x^{*}$ is close to $x_{i}$, or $z_{0}=x_{i+1}$, $z_{1}=x_{i}$ else.
The point $z_{2}$ will by the nearest point $x_{j}, j \in\{0,1, \ldots, i-1, i+2, \ldots, n\}$ to $x *$ and so on and so forth.
Thus we obtain a new set of points $z_{0}, z_{1}, \ldots, z_{n}$ which have the same values as $x_{0}, x_{1}, \ldots, x_{n}$ and which are ordered by means of distance to the point $x^{*}$ :

$$
\left|x^{*}-z_{j}\right| \geq\left|x^{*}-z_{k}\right|, \quad k=0,1, \ldots, j-1, \quad j=0,1, \ldots, n
$$

In this way we can construct, successively, the interpolating polynomials $P_{i}(x), i=$ $0,1, \ldots, n$ associated to the interpolating points $z_{0}, z_{1}, \ldots, z_{i}, i=1,2, \ldots, n$.

## 4. Practical implementation

Based on the above analysis one can derive the following polynomial interpolation algorithm for estimating $f(x)$ at a point $x=x^{*}$, with a given accuracy $\varepsilon$, using the interpolation points $x_{i}, i=\overline{0, n}$ :

1. determination of interpolating points $x_{p}, x_{p+1}$ so that $x_{p}<x^{*}<x_{p+1}$;
2. if $\left|x^{*}-x_{p}\right|>\left|x^{*}-x_{p+1}\right|$ then $\left\{\begin{array}{l}z_{0}=x_{p+1} \\ z_{1}=x_{p}\end{array}\right.$, else $\left\{\begin{array}{l}z_{0}=x_{p} \\ z_{1}=x_{p+1}\end{array}\right.$
3. calculus of $P_{1}\left(x^{*}\right), P_{1}(x)$ being the chosen interpolating polynomial with respect to the points $z_{0}, z_{1}$;
4. determination of the following interpolation point $z_{2} \in\left\{x_{0}, x_{1}, \ldots, x_{p-1}, x_{p+2}, \ldots, x_{n}\right\}$ so that

$$
\left|x^{*}-z_{2}\right|=\min \left\{\left|x^{*}-x_{i}\right|, i \in\{0,1, \ldots, p-1, p+2, \ldots, n\}\right\}
$$

5. calculus of $P_{2}\left(x^{*}\right), P_{2}(x)$ being the chosen interpolating polynomial with respect to the points $z_{0}, z_{1}, z_{2}$;
6. if $\left|P_{1}\left(x^{*}\right)-P_{2}\left(x^{*}\right)\right| \leq \varepsilon$ then $f\left(x^{*}\right) \cong P_{2}\left(x^{*}\right)$;
if not the case, we follow the steps 4-6 again, in order to add the following interpolation points $z_{i}, i=3,4, \ldots$ and to set up interpolating polynomials with respect to these new points, until the stopping criterion is accomplished.

Remark 4.1. (i). If the successive differences $\left|P_{i}\left(x^{*}\right)-P_{i-1}\left(x^{*}\right)\right|$ begin to increase as $i$ increases, one stops the algorithm concluding that $\varepsilon$ was chosen unreasonably. In this situation one may consider that $f\left(x^{*}\right) \cong P_{i-1}\left(x^{*}\right)$ represents a better approximation; (ii). In order to obtain a minimal computational cost for the above algorithm we have to take into consideration that the evaluation of the interpolating polynomials $P_{k}\left(x^{*}\right), k=1,2, \ldots$ must be done in a manner that uses the previous calculations to as greater advantage as possible, (see [1]);
(iii). The algorithm can be used for evaluate $f(x)$ at a sequence of points $x_{i}^{*}, i=$ $0,1, \ldots, r, r \in \mathrm{~N}^{*}$.

## 5. Numerical examples

5.1 We consider the data

| $x_{i}$ | 0 | 0.1 | 0.2 | 0.3 | 0.4 | 0.78 | 1.33 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $f_{i}$ | -1 | -0.620500 | -0.283987 | 0.006601 | 0.248424 | 0.677713 | -0.230627 |

We wish to approximate $f(0.155), f(0.947)$ within a given accuracy $\varepsilon$.
Using the presented algorithm we obtain:

| $x^{*}$ | accuracy <br> $\varepsilon$ | type of interpolating <br> polynomial | number of used <br> interpolating points | approximate value <br> of $f\left(x^{*}\right)$ |
| :---: | :---: | :---: | :---: | :---: |
| 0.155 | $1 \mathrm{e}-2$ | Lagrange | 3 | -0.4297343 |
|  | $1 \mathrm{e}-3$ | Lagrange | 4 | -0.4299099 |
|  | $1 \mathrm{e}-5$ | Lagrange | 5 | -0.4299075 |
|  | $2 \mathrm{e}-6$ | Lagrange | 6 | -0.4299082 |
| 0.947 | $3 \mathrm{e}-2$ | Newton | 4 | 0.6035864 |
|  | $1 \mathrm{e}-3$ | Newton | 6 | 0.6004708 |
|  | $1 \mathrm{e}-4$ | Newton | 7 | 0.6005317 |

Remark 5.1. The exact values are $\left\{\begin{array}{l}f(0.155)=-0.4299082 \\ f(0.947)=0.6005443\end{array}\right.$
5.2 We consider the data

| $x_{i}$ | 1 | 1.2 | 1.5 | 1.65 | 2.3 | 2.8 | 4.3 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $f_{i}$ | 1 | 1.648721 | 2.718282 | 3.490343 | 7.389056 | 14.154040 | 20.085537 |

We wish to approximate $f(0.3), f(2.7)$ within a given accuracy $\varepsilon$.
Using the presented algorithm we obtain:

| $x^{*}$ | accuracy <br> $\varepsilon$ | type of interpolating <br> polynomial | number of used <br> interpolating points | approximate value <br> of $f\left(x^{*}\right)$ |
| :---: | :---: | :---: | :---: | :---: |
| 0.3 | $1 \mathrm{e}-2$ | Newton | 6 | 1.3485042 |
|  | $3 \mathrm{e}-3$ | Newton | 6 | 1.3505583 |
| 2.7 | $1 \mathrm{e}-2$ | Lagrange | 7 | 14.8815651 |
|  | $4 \mathrm{e}-4$ | Lagrangee | 7 | 14.8798981 |

Remark 5.2. The exact values are $\left\{\begin{array}{l}f(0.3)=1.3498595 \\ f(2.7)=14.8797323\end{array}\right.$
5.3 We consider the data

| $x_{i}$ | 1 | 1.2 | 1.5 | 1.65 | 2.3 | 2.8 | 4.3 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $f_{i}$ | 1.684370 | 2.199796 | 2.895113 | 3.223371 | 4.579691 | 5.592577 | 8.599632 |
| $f_{i}^{\prime}$ | 2.742245 | 2.443303 | 2.221171 | 2.159282 | 2.041032 | 2.014902 | 2.000737 |

We wish to approximate $f(1.8), f(3.1)$ within a given accuracy $\varepsilon$.
Using the presented algorithm we obtain:

| $x^{*}$ | accuracy <br> $\varepsilon$ | type of interpolating <br> polynomial | number of used <br> interpolating points | approximate value <br> of $f\left(x^{*}\right)$ |
| :---: | :---: | :---: | :---: | :---: |
| 1.8 | 1e-3 | Hermite | 3 | 3.5438120 |
|  | $3 \mathrm{e}-5$ | Hermite | 4 | 3.5438051 |
|  | $4 \mathrm{e}-7$ | Hermite | 6 | 3.5438046 |
| 3.1 | $1 \mathrm{e}-3$ | Hermite | 3 | 6.1959715 |
|  | $1 \mathrm{e}-4$ | Hermite | 4 | 6.1959610 |

Remark 5.3. The exact values are $\left\{\begin{array}{l}f(1.8)=3.5438026 \\ f(3.1)=6.1959327\end{array}\right.$

## 6. Conclusions

In this paper we present an algorithm for estimating a numerical function $f(x)$ at a specific point $x=x^{*}$, based on the polynomial interpolation and the accuracy of interpolation.
Starting with an estimation of the size of the error we derive an algorithm which purpose is to evaluate $f\left(x^{*}\right)$ within a given accuracy and a minimal computational cost.
The numerical experiments performed by practical implementation of the above algorithm lead to the conclusion that this one offers a fast tool to evaluate $f(x)$ at a point $x^{*}$ with a given reasonable accuracy $\varepsilon$, by means of an interpolating polynomial of suitable degree, and in consequence with a reduced amount of work.

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(Romulus Militaru) Departamentul de Matematici Aplicate, Universitatea din Craiova Al. I. Cuza street, 13, Craiova RO-200585, Romania
E-mail address: militaruromulus@yahoo.com

