

On Metric Dimension of Hendecagonal Circular Ladder H_n

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ABSTRACT. Let $\zeta = (V, E)$ be a n th order connected graph. If the distance vectors to the vertices in an ordered subset G of vertices can uniquely identify each vertex of the graph ζ , then the set G is known as resolving set for the graph ζ . The resolving set G with smallest cardinality serves as the metric basis of graph ζ and the cardinality of this smallest resolving set serves as metric dimension for ζ . In this article, two families of convex polytopes that are closely linked are demonstrated and it is found that the metric dimension is three for both the families.

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1. Introduction

The metric dimension was found to be applicable in a variety of science and technology sectors. Mathematics and computer science both have multiple uses for the ideas of metric dimension and resolving set. Metric basis and resolving set are now basic aspects in combinatorial chemistry and molecular topology. Let $\zeta = (V, E)$ be a connected, simple and undirected graph, where V and E are the sets of vertices and edges respectively. Let the distance between the vertices m and n be denoted by $d_\zeta(m, n)$ (or simply $d(m, n)$). The maximum degree and minimum degree of ζ are denoted by $\Delta(\zeta)$ and $\delta(\zeta)$, respectively.

A vertex y is said to resolve two distinct vertices $t, g \in V(\zeta)$ if $d(y, t) \neq d(y, g)$. If each pair of vertices $t, g \in V(\zeta)$ with $t \neq g$ is resolved by some vertex $y \in F$, where F is a subset of vertices in $V(\zeta)$, then F is called as *resolving set* or *metric generator* for ζ . A resolving set with lowest cardinality is called as metric basis for graph ζ and this minimum cardinality is known as *metric dimension* of graph ζ and is represented by $\dim(\zeta)$. For an ordered subset of vertices $M = \{l_1, l_2, l_3, \dots, l_j\}$ the coordinate/ j -code/representation of a vertex k in $V(\zeta)$ is

$$d(k|F) = (d(l_1, k), d(l_2, k), d(l_3, k), \dots, d(l_j, k)).$$

In this regard, the metric generator for ζ is the set M , if $d(h|M) \neq d(u|M)$, for any two vertices $h, u \in V(\zeta)$ with $h \neq u$.

The metric generators were referred to as the locating sets by Slater[14], who also developed the concept of the metric dimension. On the other side, the identical concept of metric dimension was demonstrated by Harary and Melter[4], who referred to metric generators as resolving sets. Throughout this paper, we shall adopt resolving

sets for metric generators and locating sets. In addition to these two [4, 11] significant pioneering publications, there are other studies on applications and theoretical aspects of this invariant that may be found in the literature [1, 3, 7, 8, 9].

In basic geometry, polytope is a geometrical object with flat sides. A special case of polytopes is convex polytope which is also a polytope with the additional property of convexity and contained in the n -dimensional space \mathbb{R}^n . Applications for convex polytopes can be found throughout many branches of mathematics and computer science. Some kinds of convex polytopes' metric dimensions have been taken into consideration in [5, 6, 10]. Sharma and Bhat [11, 12, 13] introduced some plane graphs with constant metric dimension. Next, we describe some of the fundamental characteristics of a new family of convex polytopes.

Hendecagonal Circular Ladder: The hendecagonal circular ladder (HCL), designated by H_n , is a convex polytope with radial symmetry with $6n$ vertices and $7n$ edges. There are $2n$ vertices having degree 3 and $4n$ vertices having degree 2 in it. The hendecagonal circular ladder is made up of n faces with 11 sides each, one face with n sides, and another face with $2n$ sides, as shown in Figure 1.

An open problem was brought forward by Imran et al. [5] that:

Problem: *Characterize the classes of radially symmetrical plane graphs H obtained from Γ by introducing new edges in Γ such that $\dim(H) = \dim(\Gamma)$ and $V(H) = V(\Gamma)$.*

We create a plane graph family, H_n , as indicated above, in an effort to partially deal with this problem. In an attempt to partially answer this problem, we construct a plane graph family, H_n as defined above. Next, by introducing new edges to H_n at various locations, we create a new family of convex polytopes called H_n^* with a similar set of vertices. In this article, we computed the metric dimension for two classes of convex polytopes that share a common set of vertices and are closely connected.

2. Preliminaries

We discuss some fundamental ideas about the metric dimension of graphs in this section.

Independent Set [2]: A subset of vertices in a graph in which there is no pair of vertices that are adjacent is called as independent set.

Independent resolving set [2]: A subset B of vertices in a graph which is both independent and resolving is known as independent resolving set.

In this study, we focus on two convex polytopes for which we have $V(H_n) = V(H_n^*) = \{p_i, q_i, r_i, s_i, t_i, u_i : 1 \leq i \leq n\}$. The set of coordinates or metric codes for the vertices p_i, q_i, r_i, s_i, t_i and u_i is denoted by P, Q, R, S, T and U respectively, for the convex polytopes H_n and H_n^* .

For those graphs with two metric dimensions, Khuller et al.[7] introduced the subsequent result:

Theorem 2.1. Let $B \subset V(\zeta)$ be the metric basis for the connected graph ζ with cardinality two i.e., $|B| = 2$ and let $B = \{w, e\}$. Then, the following are true:

- (i) Between the vertices w and e , a shortest path P uniquely exists.
- (ii) The valencies of the vertices w and e cannot exceed 3.
- (iii) The valency of any other vertex on P cannot exceed 5.

3. Metric dimension of Hendecagonal Circular Ladder H_n

The structure of the new family of H_n is discussed in this section. We examine some of its fundamental attributes and determine its metric dimension.

The Graph of H_n :

The HCL H_n can be obtained from from the Heptagonal Circular Ladder Γ_n by placing $2n$ new vertices between the vertices p_t and q_t ($1 \leq t \leq n$) in Γ_n . It contains $7n$ edges and $6n$ vertices (see Fig.1). The HCL's H_n vertex set and edge set are separately portrayed by $V(H_n)$ and $E(H_n)$, where $V(H_n) = \{p_i, q_i, r_i, s_i, t_i, u_i : 1 \leq i \leq n\}$ and $E(H_n) = \{p_tq_t, q_tr_t, r_ts_t, s_tt_t, t_tu_t, u_tt_{t+1}, p_tp_{t+1} : 1 \leq t \leq n\}$.

We call the cycle that the vertices $\{p_i : 1 \leq i \leq n\}$ induced in graph H_n as p -cycle, the cycle generated by the vertices $\{t_i, u_i : 1 \leq i \leq n\}$ in H_n as the tu -cycle, the vertices $\{q_i : 1 \leq i \leq n\}$ in H_n as inner vertices, the vertices $\{s_i : 1 \leq i \leq n\}$ in the graph H_n as outer vertices. In the subsequent result, we look into the HCL H_n graph's metric dimension.

Theorem 3.1. $dim(H_n) = 3$, where $n \geq 21$ is odd integer.

Proof. Since, n is odd, so $n = 2w + 1$, where $w \geq 3$ is an integer. Let $F = \{p_2, p_{w+1}, p_n\} \subset V(H_n)$. To show that F is a resolving set for HCL H_n , we assign metric codes to each vertex of H_n with respect to the set F .

The metric co-ordinates for the vertices $\{p_l : 1 \leq j \leq n\}$ are given by

$$d(p_j|F) = \begin{cases} (1, w, 1), & j = 1; \\ (j - 2, w - j + 1, j), & 2 \leq j \leq w; \\ (j - 2, j - w - 1, 2w - j + 1), & w + 1 \leq j \leq w + 2; \\ (2w - j + 3, j - w - 1, 2w - j + 1), & w + 3 \leq j \leq 2w + 1. \end{cases}$$

The metric co-ordinates for the vertices $\{q_j : 1 \leq j \leq n\}$ are

$$d(q_j|F) = \begin{cases} (2, w + 1, 2), & j = 1; \\ (j - 1, w - j + 2, j + 1), & 2 \leq j \leq w; \\ (j - 1, j - w, 2w - j + 2), & w + 1 \leq j \leq w + 2; \\ (2w - j + 4, j - w, 2w - j + 2), & w + 3 \leq j \leq 2w + 1. \end{cases}$$

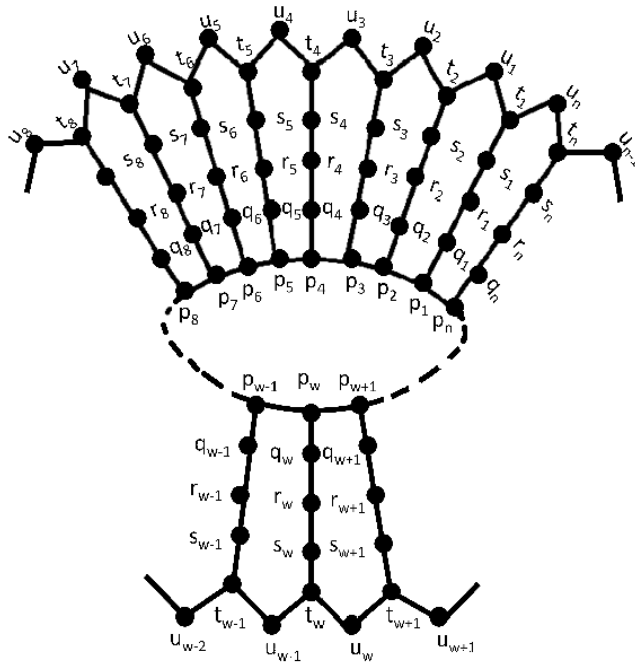


FIGURE 1. Hendecagonal Circular Ladder H_n

Next, the co-ordinates for the vertices $\{r_j : 1 \leq j \leq n\}$ are $d(r_j|F) = d(q_j|F) + (1, 1, 1)$ for $1 \leq j \leq n$, the co-ordinates for the vertices $\{s_j : 1 \leq j \leq n\}$ are $d(s_j|F) = d(q_j|F) + (2, 2, 2)$ for $1 \leq j \leq n$. Finally, the co-ordinates for the vertices of tu -cycle are $d(t_j|F) = d(q_j|F) + (3, 3, 3)$ for $1 \leq j \leq n$ and

$$d(u_j|F) = \begin{cases} (5, w + 4, 6), & j = 1; \\ (j + 3, w - j + 5, j + 5), & 2 \leq j \leq w; \\ (j + 3, j - w + 4, 2w - j + 5), & w + 1 \leq j \leq w + 2; \\ (2w - j + 7, j - w + 4, 2w - j + 5), & w + 3 \leq j \leq 2w; \\ (6, w + 5, 5), & j = 2w + 1. \end{cases}$$

From the above codes, we have $P \cap Q \cap R \cap S \cap T \cap U = \phi$, so there does not exist two vertices in H_n possessing the same metric co-ordinates, implying that $\dim(H_n) \leq 3$. Now to complete the proof, we show that $\dim(H_n) \geq 3$.

We prove that $\dim(H_n) \geq 3$ by proving that there does not exist a resolving set F with $|F| = 2$. Then, we have the following:

Resolving set	Contradictions
$F = \{p_1, p_j\}, p_j(2 \leq j \leq n)$	$d(q_1 F) = d(p_n F)$, for $2 \leq j \leq w$, and $d(u_1 F) = d(t_n F)$, for $j = w + 1$, a contradiction.
$F = \{q_1, q_j\}, q_j(2 \leq j \leq n)$	$d(q_n F) = d(p_{n-1} F)$, for $2 \leq j \leq w - 1$, $d(t_2 F) = d(r_{n-1} F)$, for $j = w$, and $d(t_2 F) = d(s_n F)$ for $j = w + 1$, a contradiction.
$F = \{r_1, r_j\}, r_j(2 \leq j \leq n)$	$d(q_n F) = d(p_{n-1} F)$, for $2 \leq j \leq w - 1$, $d(u_2 F) = d(r_n F)$, for $j = w$, and $d(u_2 F) = d(s_n F)$ for $j = w + 1$, a contradiction.
$F = \{s_1, s_j\}, s_j(2 \leq j \leq n)$	$d(q_n F) = d(p_{n-1} F)$, for $2 \leq j \leq w - 1$, $d(r_2 F) = d(p_{n-1} F)$, for $j = w$, and $d(u_2 F) = d(s_n F)$ for $j = w + 1$, a contradiction.
$F = \{t_1, t_j\}, s_j(2 \leq j \leq n)$	$d(q_{n-1} F) = d(p_{n-2} F)$, for $2 \leq j \leq w - 2$, $d(s_3 F) = d(p_n F)$, for $w - 1 \leq j \leq w$, and $d(u_2 F) = d(s_n F)$ for $j = w + 1$, a contradiction.
$F = \{u_1, u_j\}, u_j(2 \leq j \leq n)$	$d(q_{n-1} F) = d(p_{n-2} F)$, for $2 \leq j \leq w - 2$, $d(r_3 F) = d(p_1 F)$, for $w - 1 \leq j \leq w$, and $d(q_3 F) = d(p_n F)$, for $j = w + 1$, a contradiction.
$F = \{p_1, q_j\}, q_j(1 \leq j \leq n)$	$d(q_n F) = d(p_{n-1} F)$, for $1 \leq j \leq w - 1$, $d(u_1 F) = d(s_{n-1} F)$, for $j = w$, and $d(u_1 F) = d(t_n F)$, for $j = w + 1$, a contradiction.
$F = \{p_1, r_j\}, r_j(1 \leq j \leq n)$	$d(q_n F) = d(p_{n-1} F)$, for $1 \leq j \leq w - 1$, $d(u_1 F) = d(s_{n-1} F)$ for $j = w$, and $d(t_n F) = d(u_1 F)$ for $j = w + 1$, a contradiction.
$F = \{p_1, s_j\}, s_j(1 \leq j \leq n)$	$d(q_n F) = d(p_{n-1} F)$, for $1 \leq j \leq w - 1$, $d(t_1 F) = d(s_n F)$ for $j = w$, and $d(u_1 F) = d(t_n F)$ for $j = w + 1$, a contradiction.
$F = \{p_1, t_j\}, t_j(1 \leq j \leq n)$	$d(q_{n-1} F) = d(p_{n-2} F)$, for $1 \leq j \leq w - 2$, $d(q_n F) = d(p_{n-1} F)$, for $2 \leq j \leq w - 1$, $d(t_n F) = d(u_n F)$, for $j = w$, and $d(u_1 F) = d(t_n F)$ for $j = w + 1$, a contradiction.
$F = \{p_1, u_j\}, u_j(1 \leq j \leq n)$	$d(q_{n-1} F) = d(p_{n-2} F)$, for $1 \leq j \leq w - 2$, $d(q_n F) = d(p_{n-1} F)$, for $2 \leq j \leq w - 1$, $d(u_n F) = d(t_n F)$ for $j = w$, and $d(u_2 F) = d(u_{n-1} F)$ for $j = w + 1$, a contradiction.
$F = \{q_1, r_j\}, r_j(1 \leq j \leq n)$	$d(q_n F) = d(p_{n-1} F)$, for $1 \leq j \leq w - 1$, $d(t_2 F) = d(r_{n-1} F)$, for $j = w$, and $d(t_2 F) = d(s_n F)$, for $j = w + 1$, a contradiction.
$F = \{q_1, s_j\}, s_j(1 \leq j \leq n)$	$d(q_n F) = d(p_{n-1} F)$, for $1 \leq j \leq w - 1$, $d(t_2 F) = d(r_{n-1} F)$, for $j = w$, and $d(t_2 F) = d(s_n F)$, for $j = w + 1$, a contradiction.
$F = \{q_1, t_j\}, t_j(1 \leq j \leq n)$	$d(q_{n-1} F) = d(p_{n-2} F)$, for $j = 1$, $d(q_n F) = d(p_{n-1} F)$, for $2 \leq j \leq w - 1$, $d(t_2 F) = d(r_{n-1} F)$, for $j = w$, and $d(t_2 F) = d(s_n F)$, for $j = w + 1$, a contradiction.

Resolving set	Contradictions
$F = \{q_1, u_j\}, u_j(1 \leq j \leq n)$	$d(q_{n-1} F) = d(p_{n-2} F)$, for $j = 1$, $d(q_n F) = d(p_{n-1} F)$, for $2 \leq j \leq w - 1$, $d(s_2 F) = d(r_{n-1} F)$, for $j = w$, and $d(u_2 F) = d(u_{n-1} F)$, for $j = w + 1$, a contradiction.
$F = \{r_1, s_j\}, s_j(1 \leq j \leq n)$	$d(q_n F) = d(p_{n-1} F)$, for $1 \leq j \leq w - 1$, $d(s_2 F) = d(q_{n-1} F)$, for $j = w$, and $d(u_2 F) = d(s_n F)$, for $j = w + 1$ a contradiction.
$F = \{r_1, t_j\}, t_j(1 \leq j \leq n)$	$d(q_n F) = d(p_{n-1} F)$, for $2 \leq j \leq w - 1$, $d(s_2 F) = d(q_{n-1} F)$, for $j = w$, and $d(u_2 F) = d(s_n F)$, for $j = w + 1, 1$, a contradiction.
$F = \{r_1, u_j\}, u_j(1 \leq j \leq n)$	$d(q_{n-1} F) = d(p_{n-2} F)$, for $j = 1$, $d(q_n F) = d(p_{n-1} F)$, for $2 \leq j \leq w - 1$, $d(r_2 F) = d(q_{n-1} F)$, for $j = w$, and $d(u_2 F) = d(u_{n-1} F)$, for $j = w + 1$, a contradiction.
$F = \{s_1, t_j\}, t_j(1 \leq j \leq n)$	$d(q_n F) = d(p_{n-1} F)$, for $2 \leq j \leq w - 1$, $d(r_2 F) = d(p_{n-1} F)$ for $j = w$, and $d(u_2 F) = d(s_n F)$ for $j = w + 1, 1$, a contradiction.
$F = \{s_1, u_j\}, u_j(1 \leq j \leq n)$	$d(q_{n-1} F) = d(p_{n-2} F)$, for $j = 1$, $d(q_n F) = d(p_{n-1} F)$, for $2 \leq j \leq w - 1$, $d(s_3 F) = d(q_{n-1} F)$, for $j = w$, and $d(u_2 F) = d(u_{n-1} F)$, for $j = w + 1$, a contradiction.
$F = \{t_1, u_j\}, u_j(1 \leq j \leq n)$	$d(q_{n-1} F) = d(p_{n-2} F)$, for $1 \leq j \leq w - 2$, $d(p_n F) = d(s_3 F)$, for $w - 1 \leq j \leq w$, and $d(u_2 F) = d(u_{n-1} F)$, for $j = w + 1$, a contradiction.

As a result, we draw the conclusion from the explanation above that no resolving set of cardinality two exists in $V(H_n)$ implying that $dim(H_n) = 3$, which completes the proof. □

Remark 3.2. One can easily verify that the metric dimension of H_n , where n is odd with $7 \leq n \leq 19$, is also 3.

Corollary 3.3. The independent resolving number of HCL H_n is 3, where $n \geq 21$ is odd integer.

4. Metric dimension of the convex polytope H_n^*

In this section, we start by talking about the structure of a new family of H_n^* that we were able to derive from an HCL H_n . We investigate some of its fundamental properties and determine its metric dimension.

The Graph of H_n^* :

The Convex polytope H_n^* is obtained from the HCL H_n by introducing n new edges in the graph H_n between the vertices t_j and t_{j+1} for $1 \leq j \leq n$. It has $6n$ vertices and $8n$ edges. It consists of n faces having 10-sides(see Figure 2). The vertex and edge

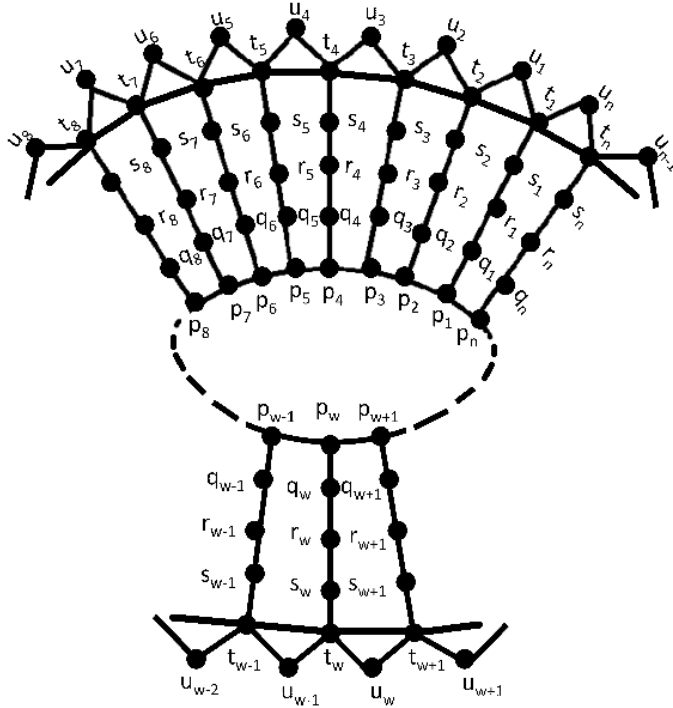


FIGURE 2. The Convex Polytope H_n^*

sets of H_n^* are depicted separately by $V(H_n^*)$ and $E(H_n^*)$, where $V(H_n^*) = V(H_n)$ and $E(H_n^*) = E(H_n) \cup \{t_j t_{j+1} : 1 \leq j \leq n\}$.

The cycle generated by the vertices set $\{p_j : 1 \leq j \leq n\}$ in the graph H_n^* as p -cycle, the cycle originated by the vertices $\{t_j, u_j : 1 \leq j \leq n\}$ in H_n^* as tu -cycle, the vertices $\{q_j : 1 \leq j \leq n\}$ in the graph H_n^* as inner vertices, the vertices $\{s_j : 1 \leq j \leq n\}$ in H_n^* as outer vertices. We examine the metric dimension of the HCL H_n^* graph in the ensuing finding.

Theorem 4.1. $dim(H_n^*) = 3$, where $n \geq 7$ is odd integer.

Proof. Since, n is odd, so $n = 2w + 1$, where $w \geq 3$ is an integer. Let $F = \{p_2, p_{w+1}, p_n\} \subset V(H_n^*)$. Now, to show that F is a resolving set for the convex polytope H_n^* , we give metric codes to each vertex of H_n^* with respect to the set F .

The metric co-ordinates for the vertices $\{p_j : 1 \leq j \leq n\}$ are

$$d(p_j|F) = \begin{cases} (1, w, 1), & j = 1; \\ (j - 2, w - j + 1, j), & 2 \leq j \leq w; \\ (j - 2, j - w - 1, 2w - j + 1), & w + 1 \leq j \leq w + 2; \\ (2w - j + 3, j - w - 1, 2w - j + 1), & w + 3 \leq j \leq 2w + 1. \end{cases}$$

The metric co-ordinates for the vertices $\{q_j : 1 \leq j \leq n\}$ are

$$d(q_j|F) = \begin{cases} (2, w + 1, 2), & j = 1; \\ (j - 1, w - j + 2, j + 1), & 2 \leq j \leq w; \\ (j - 1, j - w, 2w - j + 2), & w + 1 \leq j \leq w + 2; \\ (2w - j + 4, j - w, 2w - j + 2), & w + 3 \leq j \leq 2w + 1. \end{cases}$$

Next, the co-ordinates for the vertices $\{r_j : 1 \leq j \leq n\}$ are $d(r_j|F) = d(q_j|F) + (1, 1, 1)$ for $1 \leq j \leq n$, the co-ordinates for the vertices $\{s_j : 1 \leq j \leq n\}$ are $d(s_j|F) = d(q_j|F) + (2, 2, 2)$ for $1 \leq j \leq n$. Finally, the co-ordinates for the vertices of tu -cycle are $d(t_j|F) = d(q_j|F) + (3, 3, 3)$ for $1 \leq j \leq n$ and

$$d(u_j|F) = \begin{cases} (5, w + 4, 6), & j = 1; \\ (j + 3, w - j + 5, j + 5), & 2 \leq j \leq w; \\ (j + 3, j - w + 4, 2w - j + 5), & w + 1 \leq j \leq w + 2; \\ (2w - j + 7, j - w + 4, 2w - j + 5), & w + 3 \leq j \leq 2w; \\ (6, w + 5, 5), & j = 2w + 1. \end{cases}$$

From the above codes, we have $P \cap Q \cap R \cap S \cap T \cap U = \phi$, so we find that there does not exist a pair of vertices in H_n^* possessing the same metric co-ordinates, implying that $dim(H_n^*) \leq 3$. Now to complete the proof, it is sufficient to prove that $dim(H_n^*) \geq 3$.

We prove that $dim(H_n^*) \geq 3$ by showing that no resolving set F with $|F| = 2$ exists. Then, we have the following:

Resolving set	Contradictions
$F = \{p_1, p_j\}, p_j(2 \leq j \leq n)$	$d(q_1 F) = d(p_n F)$, for $2 \leq j \leq w$, and $d(u_1 F) = d(t_n F)$ for $j = w + 1$, a contradiction.
$F = \{q_1, q_j\}, q_j(2 \leq j \leq n)$	$d(q_n F) = d(p_{n-1} F)$, for $2 \leq j \leq w - 1$, $d(t_n F) = d(u_n F)$, for $j = w$, and $d(u_1 F) = d(t_n F)$ for $j = w + 1$, a contradiction.
$F = \{r_1, r_j\}, r_j(2 \leq j \leq n)$	$d(t_n F) = d(u_n F)$, for $2 \leq j \leq w$, $d(u_1 F) = d(t_n F)$ for $j = w + 1$, a contradiction.
$F = \{s_1, s_j\}, s_j(2 \leq j \leq n)$	$d(t_n F) = d(u_n F)$, for $2 \leq j \leq w$, and $d(u_1 F) = d(t_n F)$ for $j = w + 1$, a contradiction.
$F = \{u_1, u_j\}, u_j(2 \leq j \leq n)$	$d(t_n F) = d(u_n F)$, for $2 \leq j \leq w$, and $d(u_2 F) = d(t_n F)$ for $j = w + 1$, a contradiction.

Resolving set	Contradictions
$F = \{p_1, q_j\}, q_j(1 \leq j \leq n)$	$d(q_n F) = d(p_{n-1} F)$, for $1 \leq j \leq w - 1$, $d(t_n F) = d(u_n F)$, for $j = w$, and $d(u_1 F) = d(t_n F)$, for $j = w + 1$, a contradiction.
$F = \{p_1, r_j\}, r_j(1 \leq j \leq n)$	$d(t_n F) = d(u_n F)$, for $1 \leq j \leq w$, and $d(u_1 F) = d(t_n F)$ for $j = w + 1$, a contradiction.
$F = \{p_1, s_j\}, s_j(1 \leq j \leq n)$	$d(t_n F) = d(u_n F)$, for $1 \leq j \leq w$, and $d(u_1 F) = d(t_n F)$ for $j = w + 1$, a contradiction.
$F = \{p_1, u_j\}, u_j(1 \leq j \leq n)$	$d(t_n F) = d(u_n F)$, for $1 \leq j \leq w$, and $d(u_1 F) = d(u_n F)$ for $j = w + 1$, a contradiction.
$F = \{q_1, r_j\}, r_j(1 \leq j \leq n)$	$d(t_n F) = d(u_n F)$, for $1 \leq j \leq w$, and $d(u_1 F) = d(t_n F)$ for $j = w + 1$, a contradiction.
$F = \{q_1, s_j\}, s_j(1 \leq j \leq n)$	$d(t_n F) = d(u_n F)$, for $1 \leq j \leq w$, and $d(u_1 F) = d(t_n F)$ for $j = w + 1$, a contradiction.
$F = \{q_1, u_j\}, u_j(1 \leq j \leq n)$	$d(t_n F) = d(u_n F)$, for $1 \leq j \leq w$, and $d(u_1 F) = d(u_n F)$ for $j = w + 1$, a contradiction.
$F = \{r_1, s_j\}, s_j(1 \leq j \leq n)$	$d(t_n F) = d(u_n F)$, for $1 \leq j \leq w$, and $d(u_1 F) = d(t_n F)$ for $j = w + 1$, a contradiction.
$F = \{r_1, u_j\}, u_j(1 \leq j \leq n)$	$d(t_n F) = d(u_n F)$, for $1 \leq j \leq w$, and $d(u_1 F) = d(u_n F)$ for $j = w + 1$, a contradiction.
$F = \{s_1, u_j\}, u_j(1 \leq j \leq n)$	$d(t_n F) = d(u_n F)$, for $1 \leq j \leq w$, and $d(u_1 F) = d(u_n F)$ for $j = w + 1$, a contradiction.

As a result, we draw the conclusion from the explanation above that there does not exist resolving set of cardinality two for H_n^* implying that $dim(H_n^*) = 3$ which completes the proof. □

Corollary 4.2. The independent resolving number for the convex polytope H_n^* is 3, where $n \geq 7$ is odd integer.

5. Conclusion

This study looks at the metric dimension of an HCL H_n and a convex polytope H_n^* for odd n . We have shown for these two structures that $V(H_n) = V(H_n^*)$ and $dim(H_n) = dim(H_n^*) = 3$ for odd $n \geq 7$ (a partly solution to the question outlined in [5]). We also proved that for odd integer $n \geq 7$, the resolving set for all of these families are independent.

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